

STOCHASTIC FIXED POINTS on D

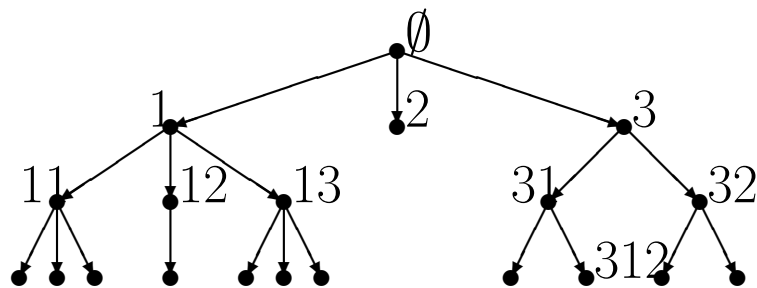
Uwe Roesler

- **Weighted Branching Process**
- **Stochastic Fixed Points on \mathbb{R}**
- **Stochastic Fixed Points on D**
- **chasing all**
- **Example Quicksort process**
If necessary
- **Quicksort Process**
- **Discrete Partial Quicksort**
- **Convergence results**

WEIGHTED BRANCHING PROCESS

Uwe Rösler

Tree V with Ulam-Harris notation



Random weights $L_e : \Omega \rightarrow G$ on edges $e = (v, vi)$,
 $(G, *)$ a semi group (with grave), often $G \subset H^H$ with
 composition

Random weights $L_e : \Omega \rightarrow G$ on pathes $e = (v, vw)$

Weighted Branching Process iff $(L_{v,vi})_{i \in \mathbb{N}}$, $v \in V$ iid

Ex: **BGW, BRW**

Ex: **Branching Processes** or **Markovian BP**

Ex: **BP with continuous time**, $L_{v,vi}$ as life time,

Ex: **Splitting processes**,

- interval splitting
- Kolmogorov rock crushing
- application in biology (genealogical tree)
- algorithms (divide and conquer algorithms, Quicksort)

Ex: **Kingman process** as time transformed splitting process

WBP, dynamic is splitting a subset $A \subset \mathbb{N}$ into two

– choose $p \in [0, 1]$ uniform distribution

– take iid Bernoulli $B_i, i \in \mathbb{N}$ to parameter p

– $A_1 := \{a \in A \mid B_a = 1\}$ $A_2 = A \setminus A_1$

– start with \mathbb{N} on root

Symmetry of edges and vertices, $V \ni v \mapsto (L_{v,vi})_i$

Skip root \emptyset , e.g. $v = (v_1, v_2, \dots, v_n) \in \mathbb{N}^*$, $L_v = L_{\emptyset, \emptyset v}$

STOCHASTIC FIXED POINTS

Stochastic fixed point equation (**SFPE**)

$$X \stackrel{\mathcal{D}}{=} f(X_1, X_2, X_3, \dots)$$

where X iid rvs and random variable f independent of X -rvs.

The distribution of X is a solution (**SFP**).

Why interested in **distribution**? Obviously exists versions of solution satisfying

$$X = f(X_1, X_2, X_3, \dots)$$

Easier and.... For simplicity on \mathbb{R} and f linear or affin

$$X \stackrel{\mathcal{D}}{=} \sum_{i \in \mathbb{N}} A_i X_i \quad X \stackrel{\mathcal{D}}{=} \sum_{i \in \mathbb{N}} A_i X_i + C$$

homogeneous sum-type and inhomogeneous sum-type.
 C stands for cost or toll.

Iterate inhomogeneous equation with equality

$$\begin{aligned} X &= C + \sum_{i \in \mathcal{N}} A_i C_i + \sum_{i, j \in \mathcal{N}} A_i A_{i, j} X_{i, j} = \dots \\ &= R_n + \sum_{|v|=n} L_v X_v \\ R_n &= \sum_{|v| < n} L_v C_v \end{aligned}$$

Choose $((L_{v, vi})_{i \in \mathcal{N}}, C_v)$, $v \in V$ as iid. Show $R_n \rightarrow_n R$ a.e. (contraction method or martingale) and $\sum_{|v|=n} L_v X_v \rightarrow 0$ disappears in distribution. Possible. Solution is R . **Endogenous** solution. **Forward** solution. Boundary at root. Probability.

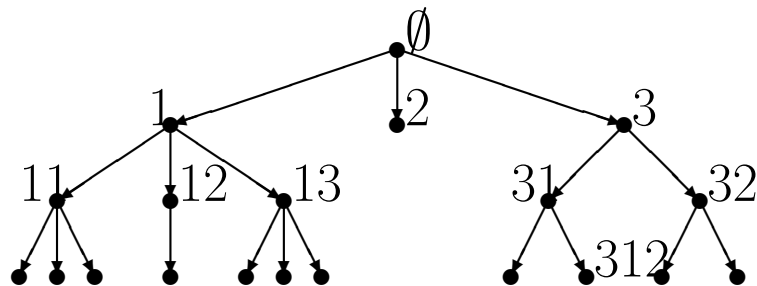
Iterate homogeneous equation with equality

$$X = \sum_{i,j \in \mathbb{N}} A_i A_{i,j} X_{i,j} = \dots = \sum_{|v|=n} L_v X_v$$

For every n exists versions of X_v , $|v| = n$ satisfying the above with equality. But in the limit???? If we consider only distributions then by Kolmogorovs projective limit exists a projective distribution. That does the job.

Backward solution. Boundary at infinity. Measure theory.

non endogenous solutions, influence of boundary



MOTIVATION for SFP

Ex: Divide and conquer algorithm provides

$$X_{n+1} \stackrel{\mathcal{D}}{=} \sum_i A_{n,i} X_{n,i} + C_n$$

assume ...

$$X \stackrel{\mathcal{D}}{=} \sum_i A_i X_i + C$$

Solve. Limiting object.

Ex: **α -stable-distributions** $\alpha \in (0, 2]$ (=Limit of sums of iid)

X α -stable distribution iff $X \stackrel{\mathcal{D}}{=} aX_1 + bX_2 + c$
for all $a, b \in \mathbb{R}$ satisfying $|a|^\alpha + |b|^\alpha = 1$ and exists c .

Replace for all by some specific a, b, c . Is solution α -stable?
Unique solution? Alsmeyer-Roesler '05

Ex: **Quicksort** Result on asymptotic running time.
Recursion for running time. Limit is SFPE.

$$X \stackrel{\mathcal{D}}{=} UX_1 + (1 - U)X_2 + C(U)$$

$$C(x) = 2x \ln x + 2(1 - x) \ln(1 - x) + 1$$

L^2 -Solution Roesler 91,92.

ALL SFPs on the REALS

Motivation for studying SFPs in its own arose by Quicksort. With contraction method invented there, 30-40 different divide and conquer algorithms could be treated, Neininger, Rüschemdorf, ... '01, '02, '04, '06 ... Metrics Wasserstein and Zolotarev

Rueschendorf: General solution of inhomogeneous is sum of general solution of homogeneous and some special solution of inhomogeneous.

Fill and Janson '00 found all solutions for Quicksort SFE, symmetric 1-stable.

Jagers and Roesler '04 Supremums and infimums type equations

SFPE on \mathbb{R}_+ and finitely many weights, Liggett, Liu on \mathbb{R} Caliebe '03, finite variance and finite first moment, (via triangle scheme)

Spitzmann: Partial results on \mathbb{R} and connection with inhomogeneous solutions

General solution on \mathbb{R} for finite number of weights, Alsmeyer-Biggins-Meiners '12 Iff conditions for all solutions. These are mixtures of α -stable distributions. α determined by $m(\alpha) = 1$ where $\beta \rightarrow m(\beta) = E(\sum_i |L_i|^\alpha)$

Solutions are

$$R + W^{1/\alpha}Y$$

$$W \stackrel{\mathcal{D}}{=} \sum_i |L_i|^\alpha W_i$$

Solutions in \mathbb{R}^d Mentemeier '12

CADLAG-PROCESSES as SFPs

$X = (X(t))_{t \in [0,1]}$ with cadlag paths

$$X \stackrel{\mathcal{D}}{=} f(X_1, X_2, X_3, \dots)$$

$f, X_i, i \in \mathbb{N}$ independent, every X_i distributed as X

Restrict us to affin f

$$X \stackrel{\mathcal{D}}{=} \sum_j T_j * X_j + C$$

where T_j is random operator on D , space and time transformation

Ex: **Brownian motion**

$$X \stackrel{\mathcal{D}}{=} \sqrt{U} X_1 + \sqrt{1-U} X_2$$

U another independent rv, values in $[0, 1]$ **or**

$$X \stackrel{\mathcal{D}}{=} (U X_1 (\frac{t}{U} \wedge 1) + (1-U) X_2 (\frac{t-U}{1-U} \vee 0))_t$$

Ex: **Cauchy**

$$X \stackrel{\mathcal{D}}{=} U X_1 + (1-U) X_2$$

solved by symmetric Cauchy(b) distribution, density $x \mapsto \frac{b}{\pi(b^2+x^2)}$ and by constants. **Or**

$$X \stackrel{\mathcal{D}}{=} (U^2 X_1 (\frac{t}{U} \wedge 1) + (1-U)^2 X_2 (\frac{t-U}{1-U} \vee 0))_t$$

Ex: **Levy processes** Analogous space-time transformation

Ex: **Find** Analysis of algorithms,
Gruebel-Roesler '96

$$X \stackrel{\mathcal{D}}{=} (\mathbb{1}_{t < U}) U X_1\left(\frac{t}{U}\right) + \mathbb{1}_{t \geq 1} (1 - U) X_2\left(\frac{t - U}{1 - U}\right) + 1$$

U uniformly distributed on $[0, 1]$

Neininger-Sulzbach '12, general functional contraction method with Zolotarev metric on D

QUICKSORT on the FLY

Conrado Martínéz: **Partial Quicksort**

Input: sequence of length n

Output: l smallest in order

Procedure: Recall Quicksort always for left most list with 2 or more elements

Publish first smallest then second smallest and so on

Observation: Algorithms does only necessary comparisons

$Y(n, l)$ number of comparisons for input $U|_n$

$$X(n, \frac{l}{n}) = \frac{Y(n, l) - EY(n, l)}{n}$$

Theo Martínéz '04

Explicit formula for $E(Y(n, l))$

Theo Roesler '13

$X(n, \dots)$ converges in Skorodhod metric almost surely to Quicksort process.

SEARCH for EXAMPLES

Quicksort, limit SFP, **All** solutions Fill-Janson '00

Quicksort process on D by contraction method

$$X \stackrel{\mathcal{D}}{=} (\mathbb{1}_{t < U} U X_1(\frac{t}{U} \wedge 1) + (1 - U) X_2(\frac{t - U}{1 - U}) + C(U, t))_t$$

U uniform independent Discrete to continuous process,
Knof '06 for finite dimensional distribution, Ragab-Roesler
'11, convergence in distribution, Roesler '13 point wise
in Skorodhod

Again endogenous and non endogenous solutions

Better work with caglad functions (caglad =left continuous with right limits)

Here functions on $[0, 1]$ and $X(0) = 0$

$$X \stackrel{\mathcal{D}}{=} (U X_1(\frac{t}{U} \wedge 1) + (1 - U) X_2(\frac{t - U}{1 - U} \vee 0) + C(U, t))_t$$

ALL CADLAG-PROCESSES as SFPs

Class of processes

$$X \stackrel{\mathcal{D}}{=} \sum_j A_j \cdot X_j \circ \varphi + C$$

A_j, C rvs D -valued, φ random time change

How many solutions has the Quicksort process SFPE?

For Quicksort: sum of ordinary QD and symmetric Cauchy distribution, Fill-Janson '00

Theo Roesler '18

For Quicksort process: sum of QP and symmetric Cauchy process

$$R + a\text{Id} + bY$$

with $a, b \in \mathbb{R}$, $b \geq 0$, R, Y independent, R solves the inhomogeneous SFPE and Y is a standard Cauchy process.

The standard Cauchy process solves the homogeneous SFPE

$$X \stackrel{\mathcal{D}}{=} (UX_1(\frac{t}{U} \wedge 1) + (1 - U)X_2(\frac{t - U}{1 - U} \vee 0))_t$$

Idea of proof: Take view of equality of random variables. and iterate via WBP

$$Y_v = \left(\mathbb{1}_{U_v > t} U Y_{v1} \left(1 \wedge \frac{t}{U_v} \right) + \mathbb{1}_{U_v \leq t} (1 - U_v) Y_{v2} \left(\frac{t - U_v}{1 - U_v} \right) \right)_t$$

Consider equations at time $t = 1$

$$Y_v(1) = U_v Y_{v1}(1) + (1 - U_v) Y_{v2}(1)$$

If we know all $Y_v(1)$ then process Y uniquely defined at the points $Y(t)$ where t are the times (places) of the pivots. These are dense in the limit and by continuity the path of Y is known.

DIRTY RECURSIONS

Recursion for Y

$$Y(x, l) = Y(x^1, l \wedge |x^1|) + Y(x^2, 0 \vee (l - |x^1| - 1)) + |x| - 1$$

and then take $x = U|_n$. Recursion for $X(n)$

Real strength of contraction method shows up for 'dirty' recursions

$$X(n) \stackrel{\mathcal{D}}{=} \sum_i L_i(I(n)) X_i(I(n)) + C(I(n))$$

$(L_i(\cdot))_i, C(\cdot), I(n), X(j), j < n$ independent, $I(n) < n$

Assume: $I(n) \rightarrow_n \infty, L_i(n) \rightarrow_n L_i, C(n) \rightarrow_n C$

Hope: $X(n) \rightarrow X$ and

$$X \stackrel{\mathcal{D}}{=} \sum_i L_i X_i + C$$

With 'nice' metric d on distributions

$d(X(n), X)$ function of $d(L_i(I(n)), L_i), d(X_i(I(n)), X_i), d(C(I(n)), C)$ shows convergence to 0.

CONVERGENCE of finite dimensional DISTRIBUTIONS

Theorem Martinez-Roesler

The one-dimensional distributions of $X(n)$ converge.

Theorem Ragab-Roesler

All finite dimensional distributions of $X(n)$ converge to the ones of the Quicksort process.

As consequence exists versions of $X(n)$ converging in Skorodhod metric to version of QP

Find algorithmic versions.

Theorem Roesler Let $U_i, i \in \mathbb{N}$, be independent uniformly distributed. Then $X(n, \cdot)$ converges almost surely to a version of the Quicksort process X in Skorodhod metric on D .

Deterministic algorithm, but random input.

SKORODHOD SPACE \mathcal{D}

\mathcal{D} equipped with Skorodhod metric d

$$d(f, g) = \inf\{\epsilon > 0 \mid \exists \lambda \in \Lambda : \|f - g \circ \lambda\|_\infty < \epsilon, \|\lambda - \text{id}\|_\infty < \epsilon\}$$

where Λ is the set of all bijective increasing functions $\lambda : [0, 1] \rightarrow [0, 1]$.

$$\begin{array}{c} \text{f} \\ \text{g} \end{array}$$

$$\begin{array}{c} \text{f} \\ \text{g} \end{array}$$

Big distance in supremum metric, small in Skorodhod metric

Theo: Alsmeyer-Biggins-Meiners

$N = \sum_i \mathbb{1}_{L_i \neq 0} < \infty$ a.e., distribution on \mathbb{R}

– $(0, \infty)$ is smallest closed multiplicative group generated by strictly positive factors

– $m(0) > 1$

– $\exists \alpha \in (0, 2] : m(\alpha) = 1$

– $\forall 0 < \beta < \alpha : 1 < m(\beta)$

Then $\psi(t) = \prod_{|v|=n} \psi_v(L_v t)$ L -a.e. for all t by martingale argument for Fourier transform. For given L

– infinitely divisible distribution,

– parameters in Levy representation satisfy fixed point equation,

– then stable.

$$\ln \psi(t) = \begin{cases} -cW|t|^\alpha(1 - i\beta \frac{t}{|t|} \tan(\frac{\pi\alpha}{2})) & \text{if } \alpha \notin \{1, 2\} \\ i\gamma Wt - cW|t| & \text{if } \alpha = 1 \\ -\sigma^2 Wt^2 & \text{if } \alpha = 2 \end{cases}$$

Remark: $N = \infty$?

Mixtures of stable distributions

INHOMOGENOUS SFE

Best result on the reals

Theo Meiners '10

Assumptions as before and on C

Set of **all** solution are distributions of

$$W^* + W^{1/\alpha}Y$$

where

- $(W^*, W), Y$ are independent,
- W^* is one solution of inhomogeneous SFE,
- $W \geq 0$ solves $W \stackrel{\mathcal{D}}{=} \sum_j |L_j|^\alpha W_j$
- Y has stable distribution to parameters

$$(\gamma, c, \beta) \in \begin{cases} \{0\} \times [-1, 1] \times [0, \infty) & \text{if } \alpha \notin \{1, 2\} \\ \mathbb{R} \times \{0\} \times [0, \infty) & \text{if } \alpha = 1 \\ \{0\} \times \{0\} \times [0, \infty) & \text{if } \alpha = 2 \end{cases}$$

SFE of SUPREMUM TYPE

On positive reals

$$X \stackrel{\mathcal{D}}{=} \sup_j L_j X_j + C$$

Jagers-Roesler, more than those expected by SFE of sum type

Rueschendorf, Alsmeyer

'complete' results